LECTURE 12

Last time:

- Strong coding theorem
- Revisiting channel and codes
- Bound on probability of error
- Error exponent

Lecture outline

- Error exponent behavior
- Expunging bad codewords
- Error exponent positivity
- Strong coding theorem

Last time

$$E_{codebooks}[P_{e,m}] \leq 2^{-N(E_0(\rho, P_X(x)) - \rho R)}$$
 for

$$E_0(\rho, P_X(x)) = -\log\left(\sum_y \left[\sum_{x_N} P_X(x) P_{Y|X}(y_i|x)^{\frac{1}{1+\rho}}\right]^{1+\rho}\right)$$

We need to:

- get rid of the expectation over codes by throwing out the worst half of the codes
- Show that the bound behaves well (exponent is $-N\alpha$ for some $\alpha > 0$)
- Relate the bound to capacity

Error exponent

Define

 $E_r(R) = \max_{0 \le \rho \le 1} \max_{P_X} (E_0(\rho, P_X(x)) - \rho R)$

then

$$E_{codebooks}[P_{e,m}] \le 2^{-NE_r(R)}$$
$$E_{codebooks,messages}[P_e] \le 2^{-NE_r(R)}$$

For a BSC:

Expunge codewords

The new $P_{e,m}$ is bounded as follows:

$$P_{e,m} = 2.2^{-NE_r(\max_{0 \le \rho \le 1} \max_{P_X} (E_0(\rho, P_X(x)) - \rho \frac{\log(2M)}{N}))} \\ = 2.2^{-NE_r(\max_{0 \le \rho \le 1} \max_{P_X} (E_0(\rho, P_X(x)) - \frac{\rho}{N} - \rho \frac{\log(M)}{N}))} \\ \le 4.2^{-NE_r(\max_{0 \le \rho \le 1} \max_{P_X} (E_0(\rho, P_X(x)) - \rho \frac{\log(M)}{N}))} \\ = 4.2^{-NE_r(R)}$$

Now we must consider positivity. Let $P_X(x)$ be such that I(X;Y) > 0, we'll show that the behavior of E_r is:

We have that:

1.
$$E_0(\rho, P_X(x)) > 0, \forall \rho > 0$$

2.
$$I(X;Y) \ge \frac{\partial E_0(\rho, P_X(x))}{\partial \rho} > 0, \forall \rho > 0$$

3.
$$\frac{\partial^2 E_0(\rho, P_X(x))}{\partial \rho^2} \le 0, \forall \rho > 0$$

We can check that

$$I(X;Y) = \frac{\partial E_0(\rho, P_X(x))}{\partial \rho}|_{\rho=0}$$

then showing 3 will establish the LHS of 2

Showing the RHS of 2 will establish 1

Let us show 3

To show concavity, we need to show that $\forall \rho_1, \rho_2 \ \forall \theta \in [0, 1]$

 $E_0(\rho_3,P_X(x))$

$$\geq \theta E_0(\rho_1, P_X(x)) + (1-\theta)E_0(\rho_2, P_X(x))$$

for
$$\rho_3 = \theta \rho_1 + \theta \rho_2$$

We shall use the fact that

$$\frac{\theta(1+\rho_1)}{1+\rho_3} + \frac{(1-\theta)(1+\rho_2)}{1+\rho_3} = 1$$

and Hölder's inequality:

$$\sum_{j} c_{j} d_{j} \leq \left(\sum_{j} c_{j}^{\frac{1}{x}}\right)^{x} \left(\sum_{j} c_{j}^{\frac{1}{1-x}}\right)^{1-x}$$

Let us pick

$$c_{j} = P_{X}(j)^{\frac{\theta(1+\rho_{3})}{1+\rho_{3}}} P_{Y|X}(k|j)^{\frac{\theta}{1+\rho_{3}}}$$

$$d_{j} = P_{X}(j)^{\frac{(1-\theta)(1+rho_{2})}{1+\rho_{3}}} P_{Y|X}(k|j)^{\frac{1-\theta}{1+\rho_{3}}}$$

$$x = \frac{\theta(1+\rho_{1})}{1+\rho_{3}}$$

Proof continued

$$\sum_{j \in \mathcal{X}} P_{X}(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_{3}}}$$

$$\leq \left(\sum_{j \in \mathcal{X}} P_{X}(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_{1}}}\right)^{\frac{\theta(1+\rho_{1})}{1+\rho_{3}}}$$

$$\times \left(\sum_{j \in \mathcal{X}} P_{X}(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_{2}}}\right)^{\frac{(1-\theta)(1+\rho_{2})}{1+\rho_{3}}}$$

$$\Rightarrow \left(\sum_{j \in \mathcal{X}} P_{X}(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_{3}}}\right)^{1+\rho_{3}}$$

$$\leq \left(\sum_{j \in \mathcal{X}} P_{X}(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_{1}}}\right)^{\theta(1+\rho_{1})}$$

$$\times \left(\sum_{j \in \mathcal{X}} P_{X}(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_{2}}}\right)^{(1-\theta)(1+\rho_{2})}$$

Proof continued

$$\Rightarrow \sum_{k \in \mathcal{Y}} \left(\sum_{j \in \mathcal{X}} P_X(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_3}} \right)^{1+\rho_3}$$

$$\le \sum_{k \in \mathcal{Y}} \left(\sum_{j \in \mathcal{X}} P_X(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_1}} \right)^{\theta(1+\rho_1)}$$

$$\times \left(\sum_{j \in \mathcal{X}} P_X(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_2}} \right)^{(1-\theta)(1+\rho_2)}$$

$$\le \left[\sum_{k \in \mathcal{Y}} \left(\sum_{j \in \mathcal{X}} P_X(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_1}} \right)^{(1+\rho_1)} \right]^{\theta}$$

$$\times \left[\left(\sum_{j \in \mathcal{X}} P_X(j) P_{Y|X}(k|j)^{\frac{1}{1+\rho_2}} \right)^{(1+\rho_2)} \right]^{(1-\theta)}$$

Proof continued

$$-\log\left(\sum_{k\in\mathcal{Y}}\left(\sum_{j\in\mathcal{X}}P_X(j)P_{Y|X}(k|j)^{\frac{1}{1+\rho_3}}\right)^{1+\rho_3}\right)$$

$$\geq -\theta\log\left(\sum_{k\in\mathcal{Y}}\left(\sum_{j\in\mathcal{X}}P_X(j)P_{Y|X}(k|j)^{\frac{1}{1+\rho_1}}\right)^{(1+\rho_1)}\right)$$

$$- (1-\theta)\left(\left(\sum_{j\in\mathcal{X}}P_X(j)P_{Y|X}(k|j)^{\frac{1}{1+\rho_2}}\right)^{(1+\rho_2)}\right)$$

$$\Rightarrow E_0(\rho_3, P_X) \geq \theta E_0(\rho_1, P_X) + (1-\theta)E_0(\rho_2, P_X)$$

so E_0 is concave!

Proof continued

Hence, extremum, if it exists, of $E_0(\rho, P_X) - \rho R$ over ρ occurs at $\frac{\partial E_0(\rho, P_X)}{\partial \rho} = R$, which implies that

$$\frac{\partial E_0(\rho, P_X)}{\partial \rho}|_{\rho=1} \le R \le \frac{\partial E_0(\rho, P_X)}{\partial \rho}|_{\rho=0} = I(X; Y)$$

is necessary for $E_r(R, P_X) = \max_{0 \le \rho \le 1} [E_0(\rho, P_X) - \rho R]$ to have a maximum

We have now placed mutual information somewhere in the expression

Critical rate is R_{CR} is defined as $\frac{\partial E_0(\rho, P_X)}{\partial \rho}|_{\rho=1}$

Proof continued

From
$$\frac{\partial E_0(\rho, P_X)}{\partial \rho} = R$$

we obtain

$$\frac{\partial R}{\partial \rho} = \frac{\partial^2 E_0(\rho, P_X)}{\partial \rho^2}$$

Hence $\frac{\partial R}{\partial \rho} < 0$, R decreases monotonically from C to R_{CR}

We can write

$$E_r(R, P_X) = E_0(\rho, P_X) - \rho \frac{\partial E_0(\rho, P_X)}{\partial \rho}$$

for $E_r(R, P_X) = E_0(\rho, P_X) - \rho R$ (ρ allows parametric relation)

then

$$\frac{\partial E_r(R, P_X)}{\partial \rho} = -\rho \frac{\partial^2 E_0(\rho, P_X)}{\partial \rho^2} > 0$$

Proof continued

Taking the ratio of the derivatives, $\frac{\partial E_r(R,P_X)}{\partial R} = -\rho$

 $E_r(R, P_X)$ is positive for R < C

Moreover

$$\frac{\partial^2 E_r(R, P_X)}{\partial R^2} = -\left[\frac{\partial^2 E_0(\rho, P_X)}{\partial \rho^2}\right]^{-1} > 0$$

thus $E_r(R, P_X)$ is convex and decreasing in R over $R_{CR} < R < C$

Proof continued

Taking $E_r(R) = \max_{P_X} E_r(R, P_X)$

is the maximum of functions that are convex and decreasing in R and so is also convex

For the P_X that yields capacity, $E_r(R, P_X)$ is positive for R < C

So we have positivity of error exponent for 0 < R < C and capacity has been introduced

This completes the coding theorem

Note: there are degenerate cases in which $\frac{\partial E_r(R,P_X)}{\partial \rho} = constant$ and $\frac{\partial^2 E_r(R,P_X)}{\partial^2 \rho} = 0$

when would that happen?

6.441 Information Theory Spring 2010

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