Clocking



Figure by MIT OCW.

Why Clocks and Storage Elements?



Digital Systems Timing Conventions

- All digital systems need a convention about when a receiver can sample an incoming data value
 - synchronous systems use a common clock
 - asynchronous systems encode "data ready" signals alongside, or encoded within, data signals
- Also need convention for when it's safe to send another value
 - synchronous systems, on next clock edge (after hold time)
 - asynchronous systems, acknowledge signal from receiver



Large Systems

 Most large scale ASICs, and systems built with these ASICs, have several synchronous clock domains connected by asynchronous communication channels



We'll focus on a single synchronous clock domain today

Clocked Storage Elements

Transparent Latch, Level Sensitive

-🛛 data passes through when clock high, latched when clock low



D-Type Register or Flip-Flop, Edge-Triggered

-🛛 data captured on rising edge of clock, held for rest of cycle



(Can also have latch transparent on clock low, or negative-edge triggered flip-flop)

Building a Latch



Static CMOS Latch Variants



node from downstream logic

Generally the best, fast and energy efficient

Weak feedback inverter so input can overpower it



Can be small, lower clock load, but sizing problematic



Latch Timing Parameters



- T_{CQmin}/T_{CQmax}
 - propagation in—out when clock opens latch
- T_{DQmin}/T_{DQmax}
 - propagation in \rightarrow out while transparent
 - usually the most important timing parameter for a latch
- T_{setup}/T_{hold}
 - define window around closing clock edge during which data must be steady to be sampled correctly

The Setup Time Race



Setup represents the race for new data to propagate around the feedback loop before clock closes the input gate.

(Here, we're rooting for the data signal)

Failing Setup



If data arrives too close to clock edge, it won't set up the feedback loop before clock closes the input transmission gate.

The Hold Time Race



Hold time represents the race for clock to close the input gate before next cycle's data disturbs the stored value. (Here we're rooting for the clock signal)

Failing Hold Time



If data changes too soon after clock edge, clock might not have had time to shut off input gate and new data will corrupt feedback loop.

Flip-Flops

 Can build a flip-flop using two latches back to back



- On positive edge, master latches input D, slave becomes transparent to pass new D to output Q
- On negative edge, slave latches current Q, master goes transparent to sample input D again

Flip-Flop Designs



- Transmission-gate master-slave latches most popular in ASICs
 - robust, convenient timing parameters, energy-efficient
- Many other ways to build a flip-flop other than transmission gate master-slave latches
 - Usually trickier timing parameters
 - -1 only found in high performance custom devices

Flip-Flop Timing Parameters



- T_{CQmin}/T_{CQmax}
 -□ propagation in→out at clock edge
- T_{setup}/T_{hold}
 - define window around rising clock edge during which data must be steady to be sampled correctly
 - either setup or hold time can be negative

Single Clock Edge-Triggered Design



Single clock with edge-triggered registers most common design style in ASICs

Slow path timing constraint

 $T_{cycle} \geq T_{CQmax} + T_{Pmax} + T_{setup} - 1 \text{ can always work around slow path by using slower clock}$

Fast path timing constraint

 $\label{eq:transform} \begin{array}{l} \mathsf{T}_{CQmin} + \mathsf{T}_{Pmin} \geq \mathsf{T}_{hold} \\ - \Box \mbox{ bad fast path cannot be fixed without redesign!} \\ - \Box \mbox{ might have to add delay into paths to satisfy hold time} \end{array}$

Clock Distribution

 Can't really distribute clock at same instant to all flip-flops on chip



Clock Grids

- One approach for low skew is to use a single metal clock grid across whole chip (Alpha 21064)
- Low skew but very high power, no clock gating



H-Trees

 Recursive pattern to distribute signals uniformly with equal delay over area



- Uses much less power than grid, but has more skew
- In practice, an approximate H-tree is used at the top level (has to route around functional blocks), with local clock buffers driving regions

Clock Oscillators

- Where does the clock signal come from?
- Simple approach: ring oscillator

Odd number of inverter stages connected in a loop

Problem:

- What frequency does the ring run at?
 - Depends on voltage, temperature, fabrication run, ...
- Where are the clock edges relative to an external observer?
 -[] Free running, no synchronization with external channel

Clock Crystals

- Fix the clock frequency by using a crystal oscillator
- Exploit peizo-electric effect in quartz to create highly resonant peak in feedback loop of oscillator
- Easy to obtain frequency accuracy of ~50 parts per million



Expensive to increase frequency to more than a few 100MHz

Phase Locked Loops (PLLs)

 Use a feedback control loop to force an oscillator to align frequency and phase with an external clock source.



Multiplying Frequency with a PLL

 By using a clock divider (a simple synchronous circuit) in the feedback loop, can force on-chip oscillator to run at rational multiple of external clock



Intel Itanium Clock Distribution



Skew Sources and Cures

- Systematic skew due to manufacturing variation can be mostly trimmed out with adaptive deskewing circuitry
 - cross chip skews of <10ps reported
- Main sources of remaining skew are temperature changes (low-frequency) and power supply noise (high frequency)
- Power supply noise affects clock buffer delay and also frequency of PLL

-1 often power for PLL is provided through separate pins

 clock buffers given large amounts of local on-chip decoupling capacitance

Skew versus Jitter

- Skew is spatial variation in clock arrival times
 - variation in when the same clock edge is seen by two different flip-flops
- Jitter is temporal variation in clock arrival times
 - variation in when two successive clock edges are seen by the same flip-flop
- Power supply noise is main source of jitter
- From now on, use "skew" as shorthand for untrimmable timing uncertainty



Skew eats into timing budget

Slow path timing constraint

 $\label{eq:tau} T_{cyc} \geq T_{CQmax} + T_{Pmax} + T_{setup} + T_{skew} \\ - \Box \mbox{ worst case is when CLK2 is earlier/later than CLK1}$

Fast path timing constraint

 $T_{CQmin} + T_{Pmin} \ge T_{hold} + T_{skew}$ -[] worst case is when CLK2 is earlier/later than CLK1